

The Path Most Travelled : Travel Demand Estimation Using Big Data Resources

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1. Algorithms

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ALGORITHM 1: Parsing OpenStreetMap Networks

```
1: {OSM files are XML based and contain way and node objects}
2: ways = set of ways in an OSM file
3: nodes = set of nodes in an OSM file
4: graph = an empty graph
5: _____
6: {Add each pair of consecutive nodes to the edge list}
7: for way in ways do
8:   for i = 0 to i = way.nodes.size() - 2 do
9:     graph.addNode(way.nodes[i])
10:    graph.addNode(way.nodes[i + 1])
11:    graph.addEdge(way.nodes[i], way.nodes[i + 1])
12: _____
13: {Simplify the network by merging road segments }
14: for way in ways do
15:   startNode = way.nodes[0]
16:   for node in way.nodes do
17:     if all edges into and out of node are segments of the same way then
18:       graph.removeNode(node)
19:       remove all edges to or from node
20:     else
21:       endNode = node
22:       graph.addEdge(startNode, endNode)
23:       FillEdgeAttributes()
24:       startNode = endNode
25: _____
26: {Notes}
27: *FillEdgeAttributes() fills in missing data such as speed limits or number
   of lanes based on way attributes
28: *graph.addNode(node) and graph.addEdge(node1, node2) only add objects
   if they do not already exist
29: *graph.removeNode(node) also removes all edges containing that node
30: *when simplifying the network, proper geographic lengths are kept even
   when nodes are deleted
```

ALGORITHM 2: Stay Point Algorithm - Step 1 - Initialize

- 1: {Each user object has a number of attributes}
- 2: *call* = a call object with an associated latitude, longitude, stay index
- 3: *calls* = vector of a user's calls ordered by timestamp
- 4: *candidateSet* = empty set of consecutive calls that meet criteria for a stay
- 5: *candidateStays* = a vector of centroids from candidate sets
- 6: δ = distance threshold between consecutive calls (in meters)
- 7: τ = time threshold between entry into and exit from the stay (in seconds)
- 8: *ds* = a grid size for the agglomerative clustering algorithm (in meters)
- 9: *stayCalls* = an empty vector of calls from stay points
- 10: {Notes}
- 11: **Centroid(callSet)* returns an object whose latitude and longitude are the centroid of all points in the input
- 12: **DistanceBetweenCalls(call1, call2)* returns the geographic distance between calls in meters
- 13: **TimeBetweenCalls(call1, call2)* returns the time between call in seconds

ALGORITHM 3: Stay Point Algorithm - Step 2 - Candidate Stays

```
1: {For each user, loop through all calls and find candidate stays}
2: candidateIndex = 0
3: candidateSet = {}
4: for  $i = 0$  to  $i = \text{calls.size}() - 2$  do
5:   if  $\text{DistanceBetweenCalls}(\text{calls}[i], \text{calls}[i + 1]) < \delta$  then
6:     candidateSet.append(calls[i + 1])
7:   else
8:     if  $\text{TimeBetweenCalls}(\text{candidateSet}[0], \text{candidateSet}[\text{end}]) > \tau$ 
       then
9:       for call in candidateSet do
10:        call.stayIndex = candidateIndex
11:        candidateStay = Centroid(candidateSet)
12:        candidateStays.append(candidateStay)
13:       candidateSet = {calls[i]}
14:       candidateIndex = candidateIndex + 1
```

ALGORITHM 4: Stay Point Algorithm - Step 3 - Agglomerative Clustering

- 1: *grid* = construct a uniform grid that covers all of a user's calls with cell dimensions $ds \times ds$
- 2: *stayIndex* = 0
- 3: **for** grid cells containing a *candidateStay* **do**
- 4: *candidateStays* = {list of candidateStay in cell}
- 5: *stay* = *Centroid(candidateStays)*
- 6: **for** *call* made from a *candidateStay* in this cell **do**
- 7: *call.longitude* = *stay.longitude*
- 8: *call.latitude* = *stay.latitude*
- 9: *call.stayIndex* = *stayIndex*
- 10: *stayCalls.append(call)*
- 11: *stayIndex* = *stayIndex* + 1

ALGORITHM 5: Stay Point Algorithm - Step 4 - Final Pass

- 1: {Final pass to add any remaining calls to the stay}
- 2: **for** $i = 0$ **to** $i = calls.size()$ **do**
- 3: **if** *call* not part of a stay and $DistanceBetweenCalls(call, stay) \leq \delta$ for any *stay* **then**
- 4: *call.longitude* = *stay.longitude*
- 5: *call.latitude* = *stay.latitude*
- 6: *call.stayIndex* = *stayIndex*
- 7: *stayCalls.append(call)*
- 8: Sort *stayCalls* by timestamp
- 9: _____

ALGORITHM 6: OD Creation Algorithm - Step 1 - Home / Work Expansion

- 1: {Data objects}
- 2: $tracts$ = census tract data objects containing demographic variables
- 3: $OD(o, d, p, t) = 0$ for origin o , destination d , purpose p , and period t
- 4: _____
- 5: {Detect home and work for all users and compute expansion factors}
- 6: **for** $user$ **in** $users$ **do**
- 7: $user.stays$ = vector of calls at stay points sorted by time
- 8: $user.home$ = index of stay point visited the most between 8pm and 7am
 on weekdays
- 9: $user.work$ = index of non-home stay point visited the most between 7am
 and 8pm on weekdays
- 10: **if** user visits work less than once per week **then**
- 11: $user.work = \text{null}$
- 12: **for** $stay$ **in** $user.stays$ **do**
- 13: $stay.label$ assigned as *home*, *work*, or *other*
- 14: $user.weekdays$ = number of weekdays a user records a stay
- 15: $user.workdays$ = number of weekdays a user records a stay at work
- 16: $tract[user.home].numUsers = tract[user.home].numUsers + 1$
- 17: **for** $tract$ **in** $tracts$ **do**
- 18: $tract.expansionFactor = tract.population/tract.numUsers$

ALGORITHM 7: OD Creation Algorithm - Step 2 - Trip Counting

```
1: {Count and expand trips}
2: for user in users do
3:   trips = empty vector to store trips taken by a user
4:   for i = 1 to i = user.stays.size() do
5:     s0 = user.stays[i - 1]
6:     s1 = user.stays[i]
7:     if s0 == S1 then
8:       continue
9:     if s0 and s1 are on the same effective day then
10:      trip = new trip from s0 to s1
11:      trip.purpose = PurposeFromLabels(s0, s1)
12:      trip.workday = true if workday for user, false otherwise
13:      trip.departure = GetConditionalDepartureTime(s0, s1)
14:      trips.append(trip)
15:     elses0 and s1 are not on the same effective day
16:      morning = create trip from home to first recorded stay
17:      night = create trip from last recorded stay to home
18:      trips.append(morning)
19:      trips.append(night)
```

```

20:   for trip in trips do
21:       o = trip.origin
22:       d = trip.destination
23:       p = trip.purpose
24:       t = trip.departure
25:       if trip.workday == true then
26:           flow = tract[user.home].expansionFactor/user.workdays
27:       else
28:           flow = tract[user.home].expansionFactor/user.weekdays
29:            $OD(o, d, p, t) = OD(o, d, p, t) + flow$ 
30: _____
31: {Notes}
32: *PurposeFromLabels(s0, s1) returns a trip purpose (HBW, NHB, HBO)
    based on the label of origin and destination stays
33: *GetConditionalDepartureTime(s0, s1) returns a departure time based on
    the observation times at origin and destination
34: *an effective day is defined as a period between 3am today until 3am on the
    next consecutive morning

```

ALGORITHM 8: Incremental Traffic Assignment

graph = road network
OD(*p*, *t*) = origin-destination matrix for purpose *p* and time window *t*
B = a bipartite network containing roads and census tracts
incrSize = vector of increment sizes, e.g. [0.4, 0.3, 0.2, 0.1]
nBatches = number of threads to use
for *i* = 0 **to** *i* < *incrSize.size()* **do**
 for *b* = 0 **to** *b* < *nBatches* **do**
 create new thread
 batch = *GetBatch*(*OD*, *b*)
 for all *o, d* pairs in *batch* **do**
 flow = *OD*[*o, d*].*flow* · *incrSize*[*i*]
 route = *A**(*o, d, graph*)
 for all segment *s* in *route* **do**
 s.flow = *s.flow* + *flow*
 $B_{e \rightarrow o} = B_{s \rightarrow o} + flow$
 wait for all threads to finish
 for segment *s* in *graph* **do**
 $s.cost \leftarrow s.freeFlowTime \cdot (1 + \alpha(\frac{s.volume}{s.capacity})^\beta)$

* *GetBatch*(*OD*, *B*) returns only the subset of *OD* pairs pertaining to a batch

* *A**(*o, d, graph*) returns the shortest path between *o* and *d* if a path exists